



ÉCOLE NATIONALE SUPÉRIEURE DES MINES

# Vehicle routing problems with road-network information

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## Vehicle Routing Problems

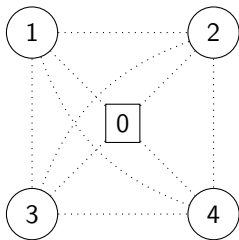
Given a complete graph  $G = (V, A)$  with  $V = \{0, \dots, n\}$

- 0 is a depot where is available a fleet of vehicles of capacity  $Q$
- nodes  $\{1, \dots, n\}$  are customers with a delivery demand  $q_i$

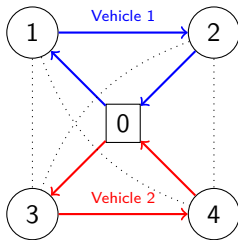
Given costs  $c_{ij}$  on arcs

Find a set of vehicle routes that serve all customers at a minimal total cost

## Vehicle Routing Problems



(a) Graph  $G$



(b) Solution

We call graph  $G$  customer-based graph

Arcs in  $G$  represent best paths in the original road-network

## THE TRUCK DISPATCHING PROBLEM\*

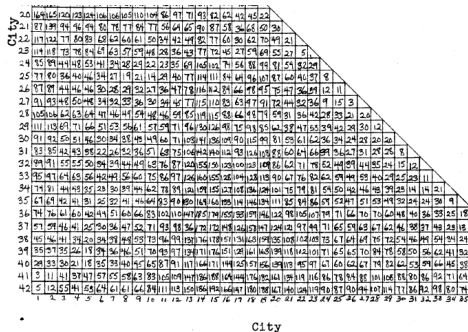
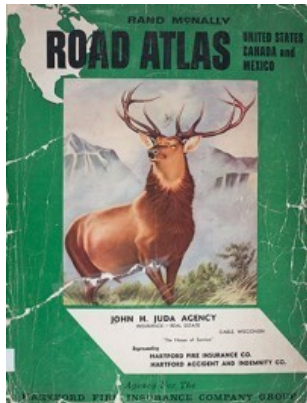
G. B. DANTZIG<sup>1</sup> AND J. H. RAMSER<sup>2</sup>

The paper is concerned with the optimum routing of a fleet of gasoline delivery trucks between a bulk terminal and a large number of service stations supplied by the terminal. The shortest routes between any two points in the system are given and a demand for one or several products is specified for a number of stations within the distribution system. It is desired to find a way to assign stations to trucks in such a manner that station demands are satisfied and total mileage covered by the fleet is a minimum. A procedure based on a linear programming formulation is given for obtaining a near optimal solution. The calculations may be readily performed by hand or by an automatic digital computing machine. No practical applications of the method have been made as yet. A number of trial problems have been calculated, however.

First paper published on the VRP

See : *G.B Dantzig, J.H. Ramser, The truck dispatching problem, Management Science, 1959*

## The first VRP

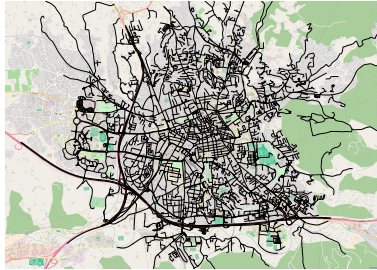


**Figure:** Geographical Information System (GIS)?  
The Rand Mc Nally road atlas (1958)

# VRPs nowadays

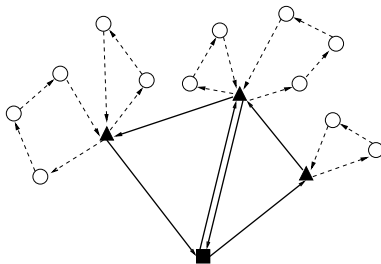


**Trend 1:** A lot of papers on urban distribution (city logistics)



## Trend 2: Accurate data

- Geographic Information Systems (openStreetMap...)
- Traffic information (historical / real-time)
- Real-time monitoring



## Trend 3: Complex organizations / models

- Time constraints
- Multiple trips
- Multiple echelons (synchronization)
- Electric vehicles (range anxiety / recharging)
- Dynamic problem...



## Outline of the presentation

### PART I: New issues

- 1 Model granularity
- 2 Complex attributes
- 3 Multiple attributes

All these issues show the limits of the customer-based graph

### PART II: Methodology

- 1 Multigraph
- 2 Road-network graph

## 1 Model granularity

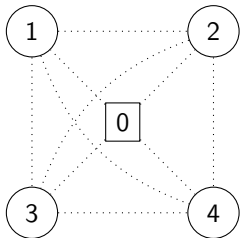
In the context of urban delivery, the distance between customers is often limited and the detail of operations (parking... ) at customers becomes important.

Typical size of a parcel delivery tour

- 40 customers
- $\leq 5$  minutes per customer, including service and traveling

See : *L. Bodin, V. Maniezzo, A. Mingozzi, Street routing and scheduling problems, in: Handbook of Transportation Science, 1999*

## 1 Model granularity



The classical model implicitly assumes:

- A unique and available parking location
- “Independence” of successive arcs in a tour

## 1 Model granularity

A unique and available parking location?

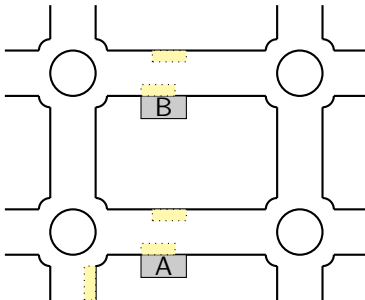
In practice :

- parking in cities is complex
- several parking locations are possible
- some booking systems start being developed

See: *Z. Lang, E. Yao, W. Hu, Z. Pan, A vehicle routing problem solution considering alternative stop points, Procedia Social and Behavioral Sciences, 2014*

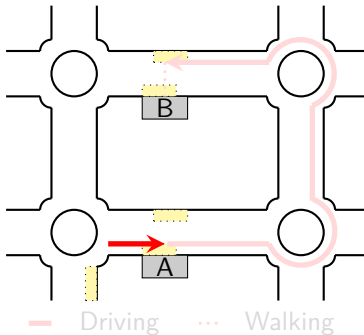
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A unique and available parking location?



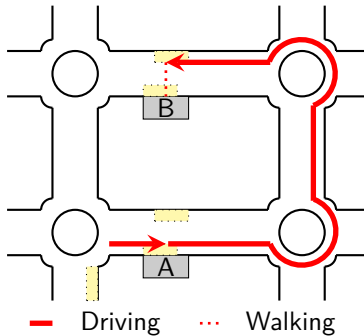
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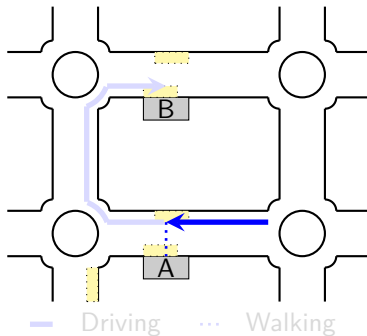
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# 1 Model granularity

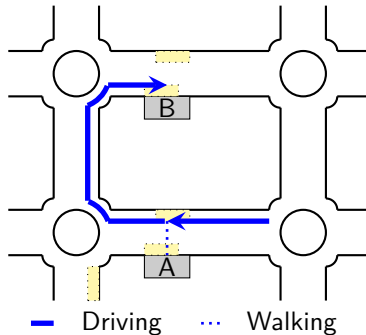
A unique and available parking location?





# 1 Model granularity

A unique and available parking location?



## 1 Model granularity

“Independence of successive arcs in a tour”?

Parking selection implies dependence between the ingoing and the outgoing arcs

This dependence also exists when some roads are subject to fees

See: *L. B. Reinhardt, M. K. Jepsen, D. Pisinger, The edge set cost of the vehicle routing problem with time windows, Transportation Science 2015.*

## 2 Complex attributes

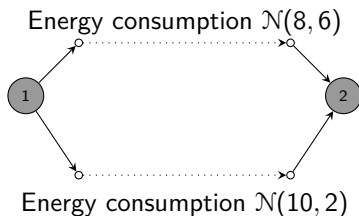
Some complications could arise with:

- Complex cost functions / constraints (fuel consumption minimization, congestion charges, etc.)
- Additional decisions
  - ▶ Breaks (driver working hour regulation)
  - ▶ Speed (speed optimization)

Not possible / not efficient (?) to precompute paths.

### Complex cost functions / constraints

**Illustration:** Electric vehicle routing with stochastic energy consumption



Best path?

## 2 Complex attributes

### Complex cost functions / constraints

**Illustration:** Electric vehicle routing with deterministic energy consumption depending on street segment slopes

Energy consumption 8 with peak at 14



Best path?

Energy consumption 10 with peak at 12

## 2 Complex attributes

Additional decisions: breaks (driver working hour regulation)

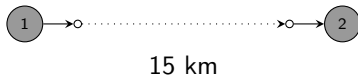


	Time intervals (min)		
	[0, 20[	[20, 40[	[40, 60]
Speed (km/h)	30	15	30

Break time: 20 minutes

## 2 Complex attributes

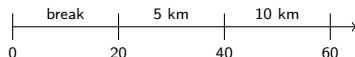
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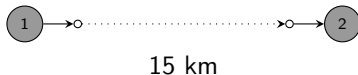
Break at customer 1:



⇒ Customer 2 reached at time 60

## 2 Complex attributes

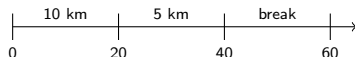
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	Time intervals (min)		
	$[0, 20[$	$[20, 40[$	$[40, 60]$
Speed (km/h)	30	15	30

Break time: 20 minutes

Break at customer 2:

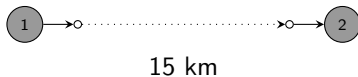


⇒ Customer 2 reached at time 40,  
break finished at time 60



## 2 Complex attributes

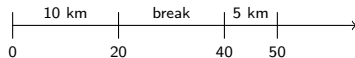
Additional decisions: breaks (driver working hour regulation)



	Time intervals (min)		
	[0, 20[	[20, 40[	[40, 60]
Speed (km/h)	30	15	30

Break time: 20 minutes

Break optimized on the road-network:



⇒ Customer 2 reached at time 50

See: *M. Chassaing, C. Duhamel, P. Lacomme. Time Dependent Capacitated Vehicle Routing Problem with Waiting Times at nodes. Odysseus, 2015.*

## 2 Complex attributes

Additional decisions: breaks (driver working hour regulation)

- In more complex networks, the path between customers 1 and 2 might even depends on the break time
- Also, the break time influences the previous / following parts of the route
- It is even possible that no solution exists with breaks at customers 1 or 2

## 2 Complex attributes

### Additional decisions: Speed

It is assumed that the decision-maker can control driver's speed to limit fuel consumption / pollution

Travel time / fuel consumption / pollution simple functions of the speed? No!



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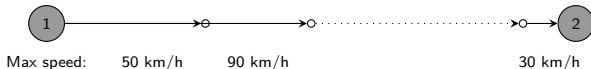


## 2 Complex attributes

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Travel time / fuel consumption / pollution simple functions of the speed? No!



## 2 Complex attributes

### Additional decisions: Speed

In addition:

- Depending on this speed different paths will be followed
- The decision-maker might modify the speed at any node in the road-network

See : *J. Qian, R. Eglese. Finding least Fuel Emission paths in a network with time-varying speeds. Networks, 2014.*

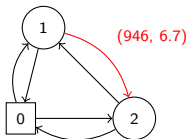
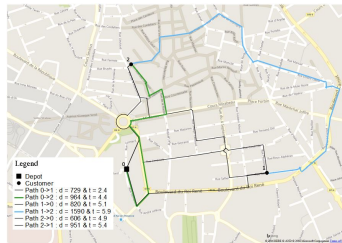
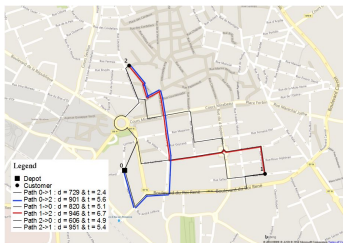
## 3 Multiple attributes

Examples of attributes:

- Distance.
- Travel time (not necessarily strongly correlated with distance).
- Energy consumption (electric vehicle), pollution, robustness, sightseeing, danger, tolls...

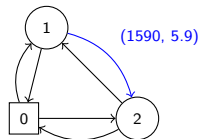
The best path is not necessarily the same for each attribute!

### 3 Multiple attributes



(a) Min-cost graph

(distance, time)



(b) Min-time graph



### 3 Multiple attributes

Some authors tried to evaluate numerically the consequences

*T. Garaix, C. Artigues, D. Feillet and D. Josselin. Vehicle routing problems with alternative paths: an application to on-demand transportation. EJOR, 2010.*

*D. Lai, O.C. Demirag and J. Leung. A tabu search heuristic for the heterogeneous vehicle routing problem on a multigraph. Transportation Research Part E, 2016*

*H. Ben-ticha, N. Absi, D. Feillet, A. Quilliot, Empirical analysis for the VRPTW with a multigraph representation for the road network, Computers & Operations Research, 2017*

Experiments show important increases of solution costs when using a customer-based graph, that can often exceed 10 %

## Outline of the presentation

### PART I: New issues

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All these issues show the limits of the customer-based graph

### PART II: Methodology

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- 2 Road-network graph

## Outline of the presentation

### Illustration with the VRP with Time Windows (VRPTW)

- Standard problem with 2 attributes: cost (distance) and time

### Methodology

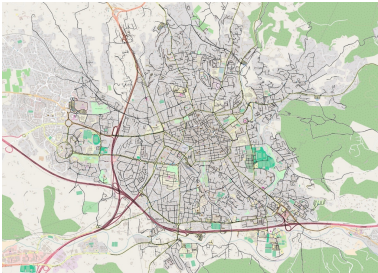
- 1 Model the road network with a multigraph.
  - ▶ One node is introduced for each customer, depot and other points of interest.
  - ▶ An arc is introduced for every efficient path between two nodes.
- 2 Apply directly solution methods on the road network.

In both cases, contrary to a customer-based graph, no information is lost.

- 1 How to construct the multigraph? Size?
- 2 How to adapt exact solution schemes
  - ▶ in multigraphs?
  - ▶ in road-network graph?
- 3 Multigraph vs road-network graph?
- 4 How to adapt heuristic solution schemes
  - ▶ in multigraphs?
  - ▶ (in road-network graphs?)

## Construction of the multigraph

Involve multi-objective shortest path problems: NP-hard



(a) 5437 nodes / 100 customers



(b) 19500 nodes / 100 customers

See: *H. Ben-ticha, N. Absi, D. Feillet, A. Quilliot, A solution method for the Multi-destination Bi-objectives Shortest Path Problem, submitted.*

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## Algorithm 1 NaiveAlgorithm(s)

---

```
1:  $\mathcal{L} \leftarrow \{(s, 0, 0)\}$  //label definition: (last vertex,distance,time)
2: repeat
3:   Select  $L = (i, d, t) \in \mathcal{L}$ 
4:    $\mathcal{L} \leftarrow \mathcal{L} \setminus \{L\}$ 
5:   for all  $j$  successor of  $i$  do
6:      $L' = (j, d + d_{ij}, t + t_{ij})$ 
7:     InsertWithDominance( $L', \mathcal{L}$ )
8:     //  $L_1 \prec L_2 \Leftrightarrow i_1 = i_2$  and  $t_1 \leq t_2$  and  $d_1 \leq d_2$ 
9:   end for
10: until  $\mathcal{L} = \emptyset$ 
```

---

Execute NaiveAlgorithm for each  $s \in V$

## Improvements:

- Implement a multi-objective multi-destination  $A^*$  to guide the search:
  - Select the label that minimizes *the detour* in distance **among all destinations**
- Stop the search once the key of the selected label is greater than the maximal detour **among all destinations**

## Other improvements with Time Windows

- Consider only *reachable* customer nodes
- Only nodes that are apt to lead to a feasible path to a destination node should be considered

## Improvements:

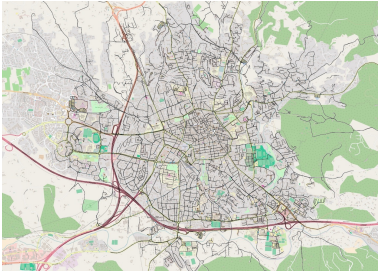
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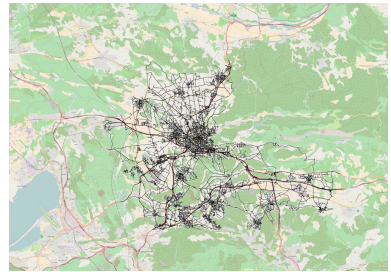
- Consider only *reachable* customer nodes
- Only nodes that are apt to lead to a feasible path to a destination node should be considered



## Construction of the multigraph



(a) 5437 nodes / 100 customers  
NaiveAlgorithm: 5 seconds  
Multi-A\*: **2 seconds**



(b) 19500 nodes / 100 customers  
NaiveAlgorithm: 400 seconds  
Multi-A\*: **30 seconds**

$\leq 4$  arcs in parallel on average

## Exact solution schemes Multigraph

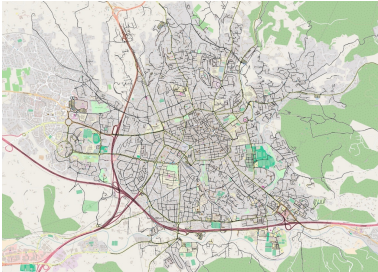
Branch-and-Price algorithms can easily be generalized

- **Master problem:** set partitioning problem
- **Pricing problem:** Elementary Shortest Path Problem with Resource Constraints on multigraph
  - Solved using an adapted labelling algorithm: a label at some node is extended to all outgoing arcs
- **Branching rules:** standard branching rules

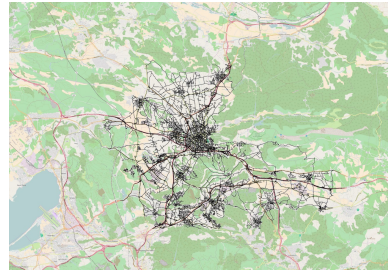
See: *H. Ben-ticha, N. Absi, D. Feillet, A. Quilliot, Empirical analysis for the VRPTW with a multigraph representation for the road network, Computers & Operations Research, 2017*

## Exact solution schemes Multigraph

Computational results: Real instances



(a) 5437 nodes / 100 customers



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# Exact solution schemes Multigraph

## Computational results: Real instances

				min-cost graph	min-time graph	multigraph	gap cost (%)	gap cost (%)
	V	C	#	CPU(s)	CPU(s)	CPU(s)	min-cost	min-time
(a)	5437	25	1	0.2	0.3	1.7	-3.4	-7.7
			2	0.2	0.2	0.8	-8.0	-6.3
		50	1	1.9	4.3	13.4	-2.3	-5.2
			2	2.9	3.7	18.8	-1.6	-4.4
		75	1	23.4	21.2	131.4	-0.5	-5.1
			2	311.6	11.4	73.4	-0.7	-5.9
(b)	19500	25	1	1.5	0.2	1.2	-6.6	-10.0
			2	0.2	0.4	1.1	-1.7	-8.7
		50	1	4.0	12.6	22.7	-0.1	-9.5
			2	10.4	13.6	13.3	-2.3	-8.3
		75	1	599.1	1372.3	174.1	-10.5	-4.1
			2	55.4	23.6	102.6	-0.9	-4.4

## Exact solution schemes Road network

### Column generation (**fractional solution**)

- The master problem is not modified
- The pricing algorithm is applied in the road-network graph
  - ▶ many nodes
  - ▶ a few arcs from each node
- The service is elementary but not the routes: crossroad nodes or arcs can be traversed many times. . .
- When extending a label to a customer, two labels are generated: one with service, one without service

See: *A. Letchford, S. Nasiri and A. Oukil. Pricing routines for vehicle routing with time windows on road networks. Computers & Operations Research, 2014.*

## Exact solution schemes Road network

### Branch-and-price (**integer solution**)

- A fractional solution can be supported by an integer flow
- No simple way to deal with it
- Same difficulties in arc routing: *C. Bode and S. Irnich. Cut-First branch-and-price-second for the capacitated arc-routing problem, Operations research 2012.*

However it “never” happens (unlike to what is happening in arc routing problems)

## Exact solution schemes Road network

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## Exact solution schemes Road network

Branch-and-price:

- When the flow is fractional: branch on arc flow
- When the flow is integer:
  - ▶ **branch 1:** enumerate all the feasible routes in the subgraph supported by the flow and solve by IP
  - ▶ **branch 2:** impose to use an arc not in the subgraph supported by the flow

$$\sum_{(i,j) \in A \setminus \tilde{A}} \sum_{r \in \Omega} b_{ijr} x_r \geq 1$$

See: *H. Ben-ticha, N. Absi, D. Feillet, A. Quilliot, T. van Woensel, A branch-and-price Algorithm for the Vehicle Routing Problem with Time Windows on a road network, submitted*



# Multigraph versus road-network Experiments

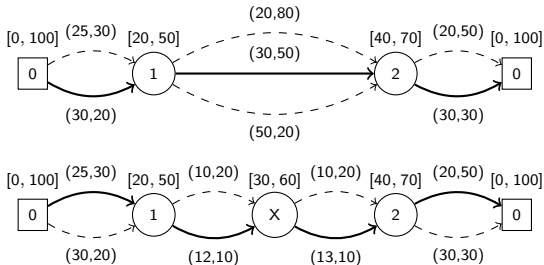
$ V_{RN} $	$ A_{RN} $	$ C $		$CPU_{MG}$	$CPU_{RN}$	$\frac{CPU_{RN}}{CPU_{MG}}$
5437	10181	5	1	0.0	5.5	117
			2	0.1	4.2	64
			3	0.0	4.3	95
			4	0.1	8.8	149
			5	0.0	14.8	330
		10	1	0.1	11.9	128
			2	0.1	7.1	90
			3	0.1	6.0	87
			4	0.1	11.6	135
			5	0.1	25.6	337
		25	1	0.2	56.9	283
			2	0.2	51.3	252
			3	0.2	35.1	213
			4	0.4	111.8	285
			5	0.2	81.0	463
		50	1	1.0	113.5	114
			2	3.3	-	-
			3	2.1	147.6	70
			4	1.0	252.0	244
			5	17.4	-	-

—: instances not solved in 7200 seconds

# Heuristic solution schemes

## Multigraph

Local search operations (e.g., an insertion, a removal) imply reoptimizing the selection of arcs.



## Heuristic solution schemes Multigraph

Arc selection is NP-hard

It can be managed by dynamic programming

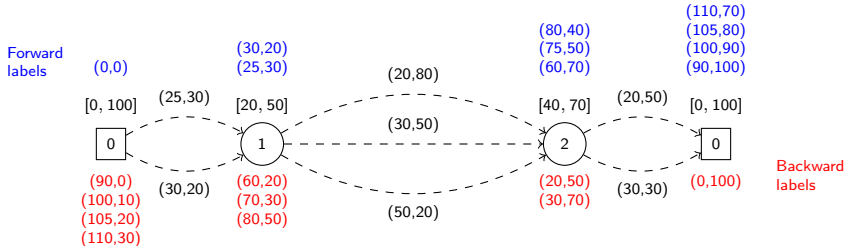
See: *T. Garaix, C. Artigues, D. Feillet and D. Josselin. Vehicle routing problems with alternative paths: an application to on-demand transportation. EJOR, 2010.*

Incremental techniques can be implemented to accelerate the method

See: *H. Ben-ticha, N. Absi, D. Feillet, A. Quilliot, T. van Woensel, Adaptive Large Neighborhood Search for VRPTW on multigraph, submitted.*

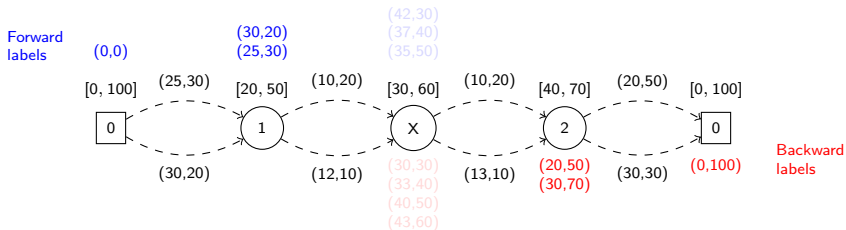
## Heuristic solution schemes Multigraph

- 1 Initially the best arcs are selected via dynamic programming (backward + forward)
- 2 Labels are stored
- 3 When a move is applied, these labels are used



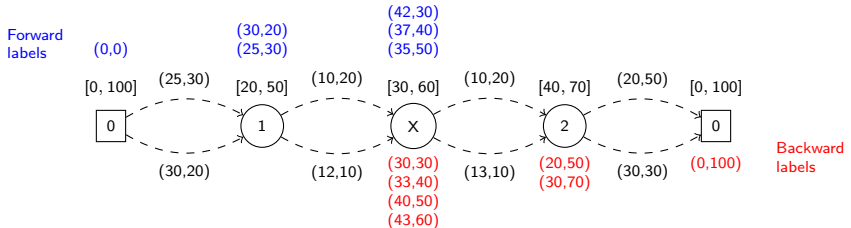
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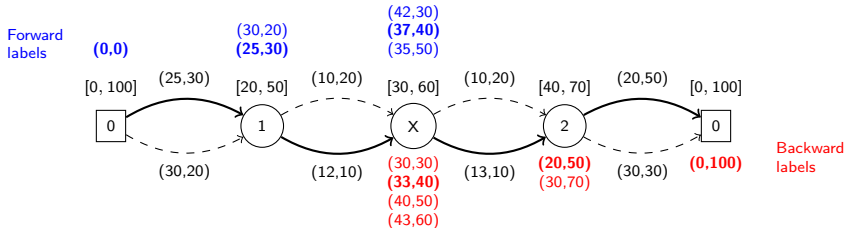
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Recent VRPs often involve

- Urban distribution
- Accurate data
- Complex organization / models

Customer-based graphs often fail modeling these VRPs with accuracy because of

- Model precision (granularity)
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The number of papers investigating these issues is very limited...  
...even if it has grown a lot recently!

Replacing the customer-based graph with a multigraph seems efficient, but is not always possible (or easy).

Replacing the customer-based graph with a road-network graph is not tractable yet.

Still a lot to do!

*H. Ben-ticha, N. Absi, D. Feillet, A. Quilliot, Vehicle routing problems with road-network information: State of the art, Networks, to appear*

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# Vehicle routing problems with road-network information

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